

THE ASYMPTOTIC BEHAVIOUR OF LOVÁSZ' ϑ FUNCTION FOR RANDOM GRAPHS

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We show that for a random graph Lovász' ϑ function is of order $\sqrt[3]{n}$.

In investigating certain properties of graphs, the behaviour of random graphs can serve as a guide. This motivates the study of the ϑ function for random graphs; one might hope to obtain a good estimate for the Shannon capacity. Unfortunately, this is not the case; the ϑ function is typically very large (of order $\sqrt[3]{n}$), while the Shannon capacity is likely to be of order $\log n$ for a random graph.

Before stating our result we recall a theorem on the eigenvalues of random matrices.

Theorem 1. (Füredi, Komlós [1]). *Let $A=(a_{ij})$ be an $n \times n$ random symmetric matrix, in which the entries a_{ij} for $i > j$ are independent identically distributed bounded random variables with distribution function H , and $a_{ii} \equiv 0$. Denote the moments of H by*

$$\mu = \int x dH(x) \quad \text{and} \quad \sigma^2 = \int x^2 dH(x) - \mu^2.$$

If $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$ are the eigenvalues of A , then

- (i) *in case $\mu > 0$ $\lambda_1 = \mu n + O(1)$ in measure, $\max_{2 \leq i \leq n} |\lambda_i| = 2\sigma \sqrt[3]{n} + O(n^{\frac{1}{3}} \log n)$ in probability,*
- (ii) *in case $\mu = 0$ $\max_{1 \leq i \leq n} |\lambda_i| = 2\sigma \sqrt[3]{n} + O(n^{\frac{1}{3}} \log n)$ in probability. ■*

A similar theorem holds for non-symmetric matrices, see [2].

Now we state our theorem on the ϑ function.

Theorem 2. *Let $\{1, \dots, n\}$ be the vertex-set of a random graph G , and denote the edge-*

set by E . The measure is $P((i, j) \in E) = p, q = 1 - p$. Then, with probability $1 - o(1)$,

$$\frac{1}{2} \sqrt{\frac{q}{p}} \sqrt{n} + O(n^{\frac{1}{3}} \log n) \leq \mathfrak{J}(G) \leq 2 \sqrt{\frac{q}{p}} \sqrt{n} + O(n^{\frac{1}{3}} \log n).$$

Proof. Define the sets \mathfrak{A} and \mathfrak{B} of symmetric matrices as follows:

$$\mathfrak{A} = \{A = (a_{ij}): a_{ij} = 1 \text{ if } (i, j) \notin E, \text{ arbitrary real otherwise}\}$$

$$\mathfrak{B} = \{B = (b_{ij}): b_{ij} = 0 \text{ if } (i, j) \in E, \text{ arbitrary real otherwise}\}.$$

It is proved in [3] that

$$\mathfrak{J}(G) = \min_{A \in \mathfrak{A}} \lambda_1(A) = \max_{B \in \mathfrak{B}} \left(1 - \frac{\lambda_1(B)}{\lambda_n(B)}\right)$$

where $\lambda_1(\cdot)$ and $\lambda_n(\cdot)$ stand for the largest and smallest eigenvalues.

Now for the upper bound in the theorem we define, for the random graph G , the matrix

$$\bar{A} = (\bar{a}_{ij}) = \begin{cases} 1 & \text{if } (i, j) \notin E \\ -\frac{q}{p} & \text{if } (i, j) \in E. \end{cases}$$

It is clear that \bar{A} is a random matrix satisfying the conditions of Theorem 1, and $\mu = 0, \sigma = \sqrt{\frac{q}{p}}$. Thus, with probability $1 - o(1)$,

$$\mathfrak{J}(G) \leq \lambda_1(\bar{A}) \leq 2\sigma \sqrt{n} + O(n^{\frac{1}{3}} \log n) = 2 \sqrt{\frac{q}{p}} \sqrt{n} + O(n^{\frac{1}{3}} \log n).$$

For proving the lower bound of our theorem, set

$$\bar{B} = (\bar{b}_{ij}) = \begin{cases} 0 & \text{if } (i, j) \in E \\ 1 & \text{if } (i, j) \notin E. \end{cases}$$

We get a random matrix with $\mu = q, \sigma = \sqrt{pq}$. Hence, with probability $1 - o(1)$,

$$\begin{aligned} \mathfrak{J}(G) &\geq 1 - \frac{\lambda_1(\bar{B})}{\lambda_n(\bar{B})} \geq \frac{\mu n + O(1)}{2\sigma \sqrt{n} + O(n^{\frac{1}{3}} \log n)} = \frac{qn + O(1)}{2 \sqrt{pq} \sqrt{n} + O(n^{\frac{1}{3}} \log n)} = \\ &= \frac{1}{2} \sqrt{\frac{q}{p}} \sqrt{n} + O(n^{\frac{1}{3}} \log n). \end{aligned}$$

Remark. One can easily extend Theorem 2 to the case $q = cn^{-1+\alpha}$. One obtains $\mathfrak{J} = O(n^{\frac{\alpha}{2}})$.

References

- [1] Z. FÜREDI and J. KOMLÓS, The eigenvalues of random symmetric matrices, *Combinatorica* **1** (1981) 233—241.
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